# A note on invariant linear transformations in multivariate public key cryptography

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## 1 Introduction

Imai and Matsumoto [1] introduced a public key cryptosystem based on multivariate quadratic polynomials. In a simplified way, the essence of their cryptosystem can be described in the following way: Start with a central monomial of the form

$$F(x) = x^{q^t + 1}$$

The secret key comprises two invertible linear transformations T and L such that

$$T\circ F\circ L$$

is the public key. In order to study equivalent public keys it is natural to ask for the "invariant" secret keys (T, L), i.e.

$$T \circ F \circ L = F$$

Lin, Faugere, Perret and Wang [2, Theorem 8] give a partial answer to this question by considering such L which fulfill

$$F \circ L = F$$

In this paper we will determine all invariant invertible linear transformations (T, L).

### 2 Preliminaries

Let K be a finite field with q elements. R is an extension field over K of degree n. We write explicitly

$$R = K[S] / \langle g(S) \rangle$$

where g(S) is an irreducible polynomial of degree n. We denote by s the image of S in R. R is a n-dimensional vector space over K. We set concretely

$$\phi: K^n \longrightarrow R$$
  $(a_0, \dots, a_{n-1}) \mapsto a_0 + a_1 s + \dots + a_{n-1} s^{n-1}$ 

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As in [2] we denote by  $\mathcal{F}$  the set of all mappings from R to R of the form

$$x \mapsto \sum_{0 \le i \le j \le n-1} \alpha_{i,j} x^{q^i + q^j}$$
 where  $\alpha_{i,j} \in R$ 

 $\mathcal{L}$  denotes the set of all invertible linear mappings from R to R of the form

$$x \mapsto \sum_{0 \le i \le n-1} \beta_i x^{q^i} \quad \text{where } \beta_i \in R$$

### 3 The Main Result

We will prove a result which extends Theorem 8 in [2]. <u>Theorem:</u> Let  $F(x) = ax^{q^i+1}$  with  $a \in R^*$  and  $1 \le i \le n-1$ . Assume we have mappings L and T in  $\mathcal{L}$  with

$$T \circ F \circ L = F$$

Then only one of the following cases can occur:

1. Case:  $4i \equiv 0 \mod n$  and  $2i \neq n$ :

$$\begin{aligned} L(x) &= cx^{q^{r}} + dx^{q^{r+2i}} \quad where \ c, d \in R^{*}, 0 \le r \le n-1 \\ T^{-1}(x) &= c^{q^{i}+1}x^{q^{r}} + cd^{q^{i}}x^{q^{r+3i}} + c^{q^{i}}dx^{q^{r+i}} + d^{q^{i}+1}x^{q^{r+2i}} \end{aligned}$$

2. Case: 2i = n:

$$L(x) = cx^{q^r}, T^{-1}(x) = (c^{q^i+1} - d)x^{q^{r+n/2}} + dx^{q^r} \qquad where \ c \in R^*, d \in R, 0 \le r \le n-1$$

3. Case:

$$L(x) = cx^{q^r}, T^{-1}(x) = c^{q^i+1}x^{q^r}$$
 where  $c \in R^*, 0 \le r \le n-1$ 

#### Notes:

1. The result in [2, Theorem 8] is covered by assuming in addition  $T^{-1}(x) = x$ . This implies in case 2

$$r = 0, d = 0, c^{q^i + 1} = 1$$

whereas in case 3 we have

$$r = 0, c^{q^i + 1} = 1$$

2. As an example for new cases compared to [2] that can occur, consider

$$F(x) = x^{q+1}, L(x) = cx + dx^{q^2}, n = 4$$

Then we have

$$F(L(x)) = c^{q+1}F(x) + cd^{q}F(x)^{q^{3}} + c^{q}dF(x)^{q} + d^{q+1}F(x)^{q^{2}}$$

3. Not every combination of c and d is possible. For instance, the condition "L is invertible" in case 1 is equivalent to the validity of the inequation

$$(-c/d)^{q^{n/2}+1} \neq 1$$

4. In case 1, T must be invertible. This is equivalent to the condition that the linear transformation

$$a_0Y + a_1Y^{q^i} + a_2Y^{q^{2i}} + a_3Y^{q^{3i}} \in R[Y]$$

in invertible where

$$a_0 = c^{q^i+1}, a_1 = c^{q^i}d, a_2 = d^{q^i+1}, a_3 = d^{q^i}c$$

We set  $p = q^{n/4}$ . p is the order of the subfield  $R_4$  of index 4 in R. If i = n/4, we can directly apply [3, p. 362] to the linear transformation

$$a_0Y + a_1Y^{q^i} + a_2Y^{q^{2i}} + a_3Y^{q^{3i}}$$
  
=  $a_0Y + a_1Y^p + a_2Y^{p^2} + a_3Y^{p^3}$ 

Therefore, T is invertible if and only if the determinant of the matrix

$$A = \begin{pmatrix} a_0 & a_3^p & a_2^{p^2} & a_1^{p^3} \\ a_1 & a_0^p & a_3^{p^2} & a_2^{p^3} \\ a_2 & a_1^p & a_0^{p^2} & a_3^{p^3} \\ a_3 & a_2^p & a_1^{p^2} & a_0^{p^3} \end{pmatrix}$$

does not vanish. In case that  $c, d \in R_4$ , this determinant is in a simple form, since then A is circular. We compute

case 
$$i = n/4; c, d \in R_4 : \det(A) \neq 0 \iff c \neq \pm d$$

If i = 3n/4, we apply [3, p. 362] to the linear transformation

$$a_0Y + a_1Y^{q^i} + a_2Y^{q^{2i}} + a_3Y^{q^{3i}}$$
  
=  $a_0Y + a_3Y^p + a_2Y^{p^2} + a_1Y^{p^3}$ 

so that we have a similar condition for the matrix A', where we exchange  $a_1$  and  $a_3$  in A. In case that  $c, d \in R_4$ , we compute the same result

case 
$$i = 3n/4; c, d \in R_4 : \det(A') \neq 0 \iff c \neq \pm d$$

<u>Proof of the main theorem</u>: It suffices to prove the result for a = 1. We extend the Frobenius mapping

$$\tau: R \longrightarrow R \qquad x \mapsto x^q$$

trivially to the polynomial ring  $R[X_0, \ldots, X_{n-1}]$ 

$$\tau: R[X_0, \dots, X_{n-1}] \longrightarrow R[X_0, \dots, X_{n-1}]$$

 $\tau$  is an isomorphism of rings which acts only on the coefficients. We set

$$Z = X_0 + X_1 s + \dots + X_{n-1} s^{n-1}$$

 $F = x^{q^i+1}$  induces via  $\phi$  a mapping from  $K^n$  to R. As this mapping, F can be written in the form of the polynomial

$$\tau^i(Z)Z \in R[X_0, \dots, X_{n-1}]$$

In the same way, the mapping

$$L(x) = \sum_{0 \le j \le n-1} \beta_j x^{q^j}$$

can be written as mapping from  $K^n$  to R in the form of the polynomial

$$\sum_{0 \le j \le n-1} \beta_j \tau^j(Z) \in R[X_0, \dots, X_{n-1}]$$

Similarly, we can write the mapping

$$T^{-1}(x) = \sum_{0 \le j \le n-1} \gamma_j x^{q^j}$$

as polynomial

$$\sum_{0 \le j \le n-1} \gamma_j \tau^j(Z) \in R[X_0, \dots, X_{n-1}]$$

The equation

$$F \circ L = T^{-1} \circ F$$

implies that the polynomial

$$\delta = \tau^i \left(\sum_{0 \le j \le n-1} \beta_j \tau^j(Z)\right) \cdot \sum_{0 \le j \le n-1} \beta_j \tau^j(Z) - \sum_{0 \le j \le n-1} \gamma_j \tau^j(\tau^i(Z)Z)$$

is - as mapping from  $K^n$  to R - identical to 0. It is well known that every  $h \in K[X_0, \ldots, X_{n-1}]$  that is identical to 0 as a mapping from  $K^n$  to K lies in the ideal

$$< X_0^q - X_0, \dots X_{n-1}^q - X_{n-1} >$$

over  $K[X_0, \ldots, X_{n-1}]$ . See for example [3, Lemma 7.40]. This immediately implies that also

$$\delta \in \langle X_0^q - X_0, \dots X_{n-1}^q - X_{n-1} \rangle$$

where the right hand side is meant as an ideal over  $R[X_0, \ldots, X_{n-1}]$ . But  $\delta$  is homogenous of degree 2. Therefore,  $\delta = 0$  as polynomial in  $R[X_0, \ldots, X_{n-1}]$  and we have the equation over  $R[X_0, \ldots, X_{n-1}]$ 

$$\tau^i(\sum_{0\leq j\leq n-1}\beta_j\tau^j(Z))\cdot\sum_{0\leq j\leq n-1}\beta_j\tau^j(Z)=\sum_{0\leq j\leq n-1}\gamma_j\tau^{i+j}(Z)\tau^j(Z)$$

 $\tau^{j}(Z)$  is the image of  $(X_0, \ldots, X_{n-1})$  under the Vandermonde matrix with entries

$$\tau^j(s^i)$$

This matrix is invertible since the elements  $\tau^{j}(s)$  are pairwise different. We define a transformation of variables by this matrix, i.e. we set

$$Y_j = \tau^j(Z)$$

in  $R[Y_0, \ldots, Y_{n-1}]$ . We write the equation above in the new variables

$$\left(\sum_{0 \le j \le n-1} \tau^i(\beta_j) Y_{i+j}\right) \cdot \sum_{0 \le j \le n-1} \beta_j Y_j = \sum_{0 \le j \le n-1} \gamma_j Y_{i+j} Y_j$$

(Indices are modulo n.) The rest of the proof follows now from the validity of this equation in an elementary, but rather tedious way: We set

$$\epsilon_t = \tau^i(\beta_{t-i})$$

so that

$$\sum_{0 \le j \le n-1} \epsilon_j Y_j = \sum_{0 \le j \le n-1} \tau^i(\beta_j) Y_{i+j}$$

We fix an index r with coefficient  $\beta_r \neq 0$ . This implies

 $\epsilon_r = 0$ 

since the monomial  $Y_r^2$  does not appear on the right side of the equation. We write the left side of the equation as sum of monomials with certain cofficients. The monomials of the form  $Y_r Y_j$  with  $j \neq r$ ,  $0 \leq j \leq n-1$ , in this sum have the coefficient

 $\epsilon_j \beta_r$ 

Because of the structure of the right side of the equation, we get  $\epsilon_j = 0$  unless j = r + ior j = r - i. This implies that in the sum

$$\sum_{0 \le j \le n-1} \beta_j Y_j$$

at most two coefficients do not vanish. Therefore, we assume that we have

$$\beta_r \neq 0$$
 and  $\beta_s \neq 0$ 

for  $r \neq s$ . The left side of the equation above now reads

$$(\tau^{i}(\beta_{r})Y_{r+i} + \tau^{i}(\beta_{s})Y_{s+i})(\beta_{r}Y_{r} + \beta_{s}Y_{s})$$
  
=  $\tau^{i}(\beta_{r})Y_{r+i}\beta_{r}Y_{r} + \tau^{i}(\beta_{r})Y_{r+i}\beta_{s}Y_{s} + \tau^{i}(\beta_{s})Y_{s+i}\beta_{r}Y_{r} + \tau^{i}(\beta_{s})Y_{s+i}\beta_{s}Y_{s}$ 

Considering the structure of the right side of the equation, this immediately implies

$$r+i-s \equiv -i \mod n \text{ and } s+i-r \equiv -i \mod n$$

This gives the condition

 $4i \equiv 0 \mod n$  and  $s \equiv r + 2i \mod n$ 

and in addition  $2i \neq n$ . The left side of the equation above now reads

$$\tau^{i}(\beta_{r})\beta_{r}Y_{r+i}Y_{r}+\tau^{i}(\beta_{r})\beta_{s}Y_{r+i}Y_{r+2i}+\tau^{i}(\beta_{s})\beta_{r}Y_{r+3i}Y_{r}+\tau^{i}(\beta_{s})\beta_{s}Y_{r+3i}Y_{r+2i}$$

Therefore, the  $\gamma_j$  are uniquely defined and the claim of case 1 follows. If in the sum

$$\sum_{0 \le j \le n-1} \beta_j Y_j$$

exactly one coefficient does not vanish, then the equation above reads

$$\tau^{i}(\beta_{r})\beta_{r}Y_{r+i}Y_{r} = \sum_{0 \le j \le n-1} \gamma_{j}Y_{i+j}Y_{j}$$

Let us assume that  $2i \neq n$ . Then, all the monomials in the sum on the right side of this equation are different. This implies directly the claim of case 3. If i = n/2, we get the equation

$$\tau^{i}(\beta_{r})\beta_{r}Y_{r+n/2}Y_{r} = \sum_{0 \le j \le n-1} \gamma_{j}Y_{n/2+j}Y_{j} = \gamma_{r}Y_{r+n/2}Y_{r} + \gamma_{r+n/2}Y_{r+n$$

which gives the claim of case 2.

#### 4 References

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[3] Lidl, Niederreiter: Finite Fields, Vol. 20, Encyclopdia of Math. and its appl., Cambridge University Press, Cambridge 1997